

#### **Data-Intensive Distributed Computing**

#### CS 431/631 451/651 (Winter 2019)

#### Part 8: Analyzing Graphs, Redux (1/2) March 21, 2019

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These slides are available at http://roegiest.com/bigdata-2019w/



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# Graph Algorithms, again? (srsly?)



Irregular structure Fun with data structures!

Irregular data access patterns Fun with architectures!

> Iterations Fun with optimizations!

## **Characteristics of Graph Algorithms**



Local computations Message passing along graph edges

Iterations

#### Visualizing Parallel BFS



#### PageRank: Defined

Given page x with inlinks  $t_1...t_n$ , where

C(t) is the out-degree of t

lpha is probability of random jump

N is the total number of nodes in the graph



#### PageRank in MapReduce



#### PageRank vs. BFS



# **Characteristics of Graph Algorithms**

#### Parallel graph traversals

Local computations Message passing along graph edges







#### MapReduce Sucks

Hadoop task startup time Stragglers Needless graph shuffling Checkpointing at each iteration

## Let's Spark!













#### MapReduce vs. Spark



Source: http://ampcamp.berkeley.edu/wp-content/uploads/2012/06/matei-zaharia-part-2-amp-camp-2012-standalone-programs.pdf

# **Characteristics of Graph Algorithms**

#### Parallel graph traversals

#### Local computations Message passing along graph edges

Iterations

Even faster?

## Big Data Processing in a Nutshell



Replicate

Reduce cross-partition communication

## Simple Partitioning Techniques

Hash partitioning

Range partitioning on some underlying linearization Web pages: lexicographic sort of domain-reversed URLs



PageRank over webgraph (40m vertices, 1.4b edges)



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# Schimmy Design Pattern

#### Basic implementation contains two dataflows: Messages (actual computations) Graph structure ("bookkeeping")

Schimmy: separate the two dataflows, shuffle only the messages Basic idea: merge join between graph structure and messages

both relationshowstations join kistently partitioned and sorted by join key







PageRank over webgraph (40m vertices, 1.4b edges)



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#### Simple Partitioning Techniques

Hash partitioning

Range partitioning on some underlying linearization Web pages: lexicographic sort of domain-reversed URLs Social networks: sort by demographic characteristics

## **Country Structure in Facebook**



Analysis of 721 million active users (May 2011)

54 countries w/ >1m active users, >50% penetration

Ugander et al. (2011) The Anatomy of the Facebook Social Graph.

### Simple Partitioning Techniques

Hash partitioning

Range partitioning on some underlying linearization Web pages: lexicographic sort of domain-reversed URLs Social networks: sort by demographic characteristics Geo data: space-filling curves



# Geo-data = regular graph



## Space-filling curves: Z-Order Curves


# Space-filling curves: Hilbert Curves



# Simple Partitioning Techniques

Hash partitioning

Range partitioning on some underlying linearization Web pages: lexicographic sort of domain-reversed URLs Social networks: sort by demographic characteristics Geo data: space-filling curves

But what about graphs in general?

Source: http://www.flickr.com/photos/fusedforces/4324320625/

## **General-Purpose Graph Partitioning**

Graph coarsening Recursive bisection

## **General-Purpose Graph Partitioning**



Karypis and Kumar. (1998) A Fast and High Quality Multilevel Scheme for Partitioning Irregular Graphs.

# Graph Coarsening



Karypis and Kumar. (1998) A Fast and High Quality Multilevel Scheme for Partitioning Irregular Graphs.

# Chicken-and-Egg

To coarsen the graph you need to identify dense local regions To identify dense local regions quickly you to need traverse local edges But to traverse local edges efficiently you need the local structure!

To efficiently partition the graph, you need to already know what the partitions are! Industry solution?

# Big Data Processing in a Nutshell

Partition

Replicate

Reduce cross-partition communication



## Partition



## Partition



#### What's the fundamental issue?

# **Characteristics of Graph Algorithms**

Parallel graph traversals

Local computations Message passing along graph edges

Iterations

## Partition



# State-of-the-Art Distributed Graph Algorithms



# Graph Processing Frameworks

Source: Wikipedia (Waste container



# **Pregel: Computational Model**

Based on Bulk Synchronous Parallel (BSP) Computational units encoded in a directed graph Computation proceeds in a series of supersteps Message passing architecture

Each vertex, at each superstep:

Receives messages directed at it from previous superstep Executes a user-defined function (modifying state) Emits messages to other vertices (for the next superstep)

#### Termination:

A vertex can choose to deactivate itself Is "woken up" if new messages received Computation halts when all vertices are inactive



Source: Malewicz et al. (2010) Pregel: A System for Large-Scale Graph Processing. SIGMOD.

# **Pregel: Implementation**

Master-Worker architecture

Vertices are hash partitioned (by default) and assigned to workers Everything happens in memory

Processing cycle:

Master tells all workers to advance a single superstep Worker delivers messages from previous superstep, executing vertex computation Messages sent asynchronously (in batches) Worker notifies master of number of active vertices

Fault tolerance

Checkpointing Heartbeat/revert

# Pregel: SSSP

```
class ShortestPathVertex : public Vertex<int, int, int> {
 void Compute(MessageIterator* msgs) {
  int mindist = IsSource(vertex id()) ? 0 : INF;
  for (; !msgs->Done(); msgs->Next())
   mindist = min(mindist, msgs->Value());
  if (mindist < GetValue()) {</pre>
   *MutableValue() = mindist;
   OutEdgeIterator iter = GetOutEdgeIterator();
   for (; !iter.Done(); iter.Next())
    SendMessageTo(iter.Target(),
            mindist + iter.GetValue());
  VoteToHalt();
 }
```

```
};
```

# Pregel: PageRank

class PageRankVertex : public Vertex<double, void, double> {
 public:

```
virtual void Compute(MessageIterator* msgs) {
```

```
if (superstep() >= 1) {
   double sum = 0;
   for (; !msgs->Done(); msgs->Next())
     sum += msgs->Value();
   *MutableValue() = 0.15 / NumVertices() + 0.85 * sum;
}
```

```
if (superstep() < 30) {
    const int64 n = GetOutEdgeIterator().size();
    SendMessageToAllNeighbors(GetValue() / n);
    } else {
    VoteToHalt();
    }
};</pre>
```

# **Pregel: Combiners**

class MinIntCombiner : public Combiner<int> {
 virtual void Combine(MessageIterator\* msgs) {

```
int mindist = INF;
for (; !msgs->Done(); msgs->Next())
mindist = min(mindist, msgs->Value());
Output("combined_source", mindist);
}
```

};

Source: Malewicz et al. (2010) Pregel: A System for Large-Scale Graph Processing. SIGMOD.



# **Giraph Architecture**

Master – Application coordinator Synchronizes supersteps Assigns partitions to workers before superstep begins

#### Workers – Computation & messaging

Handle I/O – reading and writing the graph Computation/messaging of assigned partitions

#### ZooKeeper

Maintains global application state

## **Giraph Dataflow**



# **Giraph Lifecycle**

#### Vertex Lifecycle



# **Giraph Lifecycle**



# Giraph Example

```
public class MaxComputation extends BasicComputation<IntWritable, IntWritable,</pre>
   NullWritable, IntWritable> {
 @Override
 public void compute(Vertex<IntWritable, IntWritable, NullWritable> vertex,
      Iterable<IntWritable> messages) throws IOException
  ł
    boolean changed = false;
    for (IntWritable message : messages) {
      if (vertex.getValue().get() < message.get()) {</pre>
        vertex.setValue(message);
        changed = true;
    if (getSuperstep() == 0 || changed) {
      sendMessageToAllEdges(vertex, vertex.getValue());
    vertex.voteToHalt();
```

### **Execution Trace**



Time



# State-of-the-Art Distributed Graph Algorithms



# Graph Processing Frameworks

Source: Wikipedia (Waste container

## GraphX: Motivation



# GraphX = Spark for Graphs

Integration of record-oriented and graph-oriented processing

Extends RDDs to Resilient Distributed Property Graphs

class Graph[VD, ED] {
 val vertices: VertexRDD[VD]
 val edges: EdgeRDD[ED]
}

# Property Graph: Example

#### Property Graph



#### Vertex Table

ld	Property (V)	
3	(rxin, student)	
7	(jgonzal, postdoc)	
5	(franklin, professor)	
2	(istoica, professor)	

# Edge Table

Srcld	Dstld	Property (E)
3	7	Collaborator
5	3	Advisor
2	5	Colleague
5	7	PI

# **Underneath the Covers**



# **GraphX** Operators

#### "collection" view

val vertices: VertexRDD[VD]
val edges: EdgeRDD[ED]
val triplets: RDD[EdgeTriplet[VD, ED]]

#### Transform vertices and edges

mapVertices mapEdges mapTriplets

#### Join vertices with external table

Aggregate messages within local neighborhood Pregel programs


